

Litland's "Grounding and defining identity"

1. Claims

$(\text{Null}_=) \Box_= \forall x \forall y (x = y \rightarrow ((0 \ll x = y) \wedge \forall pp (pp \ll x = y \rightarrow 0 \equiv pp)))$

- "It is essential to identity that the only ground of identity facts is 0."

$(\text{Null}_{\neq}) \Box_= \forall x \forall y (x \neq y \rightarrow ((0 \ll x \neq y) \wedge \forall pp (pp \ll x \neq y \rightarrow 0 \equiv pp)))$

- "It is essential to identity that the only ground of distinctness facts is 0."

- (i) Litland is operating with a notion of strict, full, immediate ground.
- (ii) Litland holds a parallel thesis for the identity relation of each type $\langle \tau, \tau \rangle$.
- (iii) Litland allows the essentialist operator to be indexed with items of arbitrary types.

2. Argumentative strategy

Step 1: there is a relation R such that (Null_R) and $(\text{Null}_{\lambda xy. \neg Rxy})$ follows from the real definition of R.

Step 2: R is the best candidate of =.

3. Real definition

A rough idea: to define a relation, one specifies the grounds of its instantiations and non-instantiations.

3.1. Grammar

Litland works in the language of simple, relational type theory, which contains types for pluralities.

In general, a claim of (full, immediate) real definition can be expressed as

$$s \frac{x_0 : \psi_0, x_1 : \psi_1, \dots}{x} \parallel \phi_0, \phi_1, \dots$$

which reads: s is, by definition the x such that for all x_0 that satisfies ψ_0 , x_1 that satisfies ψ_1, \dots , $\phi_0(x, x_0, x_1, \dots)$, $\phi_1(x, x_0, x_1, \dots), \dots$

Litland includes as a primitive the application relation $\mathcal{A}(Rxypp)$, which reads “ p is a result of applying R to x and y ”.

Here, the “is by definition” connective plays a complex grammatical role: it binds and generalizes variables, and restricts the range of variables it generalizes.

This complex grammar is mainly motivated by the idea that *logical operations* should be *pure*. Litland characterizes the intuitive notion of purity as

- A definition is immediately (\neg -)pure if the only items figuring in it are the grounding relation and the application relation (and negation).
- An item is immediately (\neg -)pure if its definition is immediately (\neg -)pure.
- An item is (\neg -)pure if it is either immediately (\neg -)pure or has a definition from (\neg -)pure items.

To satisfy the requirement that identity be (\neg -)pure, its definition can’t have quantifiers, conditionals, etc. in it, hence the complex grammar of the definition connective.

“Wittgensteinian variable convention”: distinct bound variables take distinct values.

3.2. Logic

Factivity: if s is by definition the x such that $\phi(x)$, then $\phi(s/x)$.

Individuation: no two items have the same real definition.

Definitional \ll -reflection: if p is, by definition, (anti-)grounded in qq , then qq are the only (immediate non-factive) (anti-)grounds of p .

Definitional \mathcal{A} -reflection: if p is, by definition, a result of applying R to x and y , then p can only results from applying R to x and y .

3.3. Metalogic

- A plurality pp of propositions is *closed* iff
 - (i) pp contains the definition of every item figuring in any proposition in p ;
 - (ii) pp is closed under Factivity, Individuation, \ll -reflection and \mathcal{A} -reflection;
 - (iii) pp contains every proposition that is grounded in some propositions amongst the pp ; and
 - (iv) any proposition amongst the pp that has some full grounds, has some full grounds amongst the pp .
- pp is *coherently closed* if it is closed doesn’t contain a contradiction.
- q is a *consequence* of pp if q is in every closed plurality containing pp .
- q is a *coherent consequence* of pp if q is in every coherently closed plurality containing pp .

- q is a *essential coherent* consequence of pp if q is a coherent consequence of pp , and every item that figures in q figures in a definition of an item in pp .

4. Really defining identity

$$= \frac{x,y,(p:\mathcal{A}(Rxxp)),(q:\mathcal{A}(Rxyq))}{R} \parallel \left\{ \begin{array}{l} p \text{ } \overline{r} \parallel \left\{ \begin{array}{l} \mathcal{A}(Rxxr) \\ 0 \ll r \end{array} \right. \\ q \text{ } \overline{r} \parallel \left\{ \begin{array}{l} \mathcal{A}(Rxyr) \\ \mathcal{A}(Ryxr) \\ 0 \ll \neg r \end{array} \right. \end{array} \right.$$

which roughly reads: identity is, by definition, the relation R such that for all x and y , for all p resulting from applying R to x and x , and for all q resulting from applying R to x and y ,

- (i) p is, by definition, a result of applying R to x and x , and is 0-grounded;
- (ii) q is, by definition, a result of applying R to x and y , and is 0-grounded.

Leibniz's Law, Strict Symmetry, and the claim that identity and distinctness facts are uniquely 0-grounded, follows from this definition and the logic sketched in 3.2.

Litland also sketches an argument to the effect that

$$\forall x \forall y (x = y \rightarrow ((0 \ll x = y) \wedge \forall pp (pp \ll x = y \rightarrow 0 \equiv pp)))$$

is an essential consequence of the real definitions of the logical operations involved. Presumably, given certain principles connecting claims involving \Box_x and claims about essential consequence that he doesn't specify, one could deduce (Null \equiv) and (Null \neq).

5. Theoretical virtues

Litland carries out Step 2 of his argumentative strategy by arguing (1) that his definition of identity validates standard logical principles involving identity; and (2) that his account has several theoretical virtues that his competitors lack.

In particular, he argues that his account

- (i) is uniform and topic-neutral;
- (ii) avoids problems that competitors have;
- (iii) avoids the objection from differential grounding,